

IBM Systems & Technology Group

z/VM 5.4 Performance Update

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z/VM Performance Evaluation

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VM Performance Evaluation

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Agenda

z/VM performance update

- z/VM 5.4 major items
- z/VM 5.4 minor items
- Interesting APARs
- Hardware notes
- z10 performance



z/VM 5.4 – Enhancements and Notes

Enhancements

- Specialty engines
- Dynamic memory upgrade
- Virtual CPU share redistribution
- DCSS above 2 GB
- TCP/IP layer 2
- Telnet IPv6
- Linux install from HMC
- Upper DAT tables
- Other changes
- Service: VMRM "safety net"
- Service: MDC

Notes

- Long eligible list stays
- Reorder processing
- VMDUMP



z/VM 5.4: Specialty Engines

- Now allow environment with real IFLs and real CPs in same LPAR (as well as other specialty engines)
 - Called "VM Mode LPAR" be careful this phrase can be ambiguous at times
- New options to SET SHARE allow a value for each processor type in a mixed virtual machine.
- CPUAFFINITY OFF will result in virtual specialty engines being dispatched on CPs
- Merging CP LPAR with IFL LPAR requires thought:
 - First step, make virtual machines on IFL LPAR have virtual IFLs
 - For duplicated work (RACF, TCP/IP, etc.), need to determine which to use or in some cases which to duplicate
 - Remember that in some environments, the IFLs may be faster than the CPs.
 - You may need to revisit your charge back model.
- See the z/VM Performance Report

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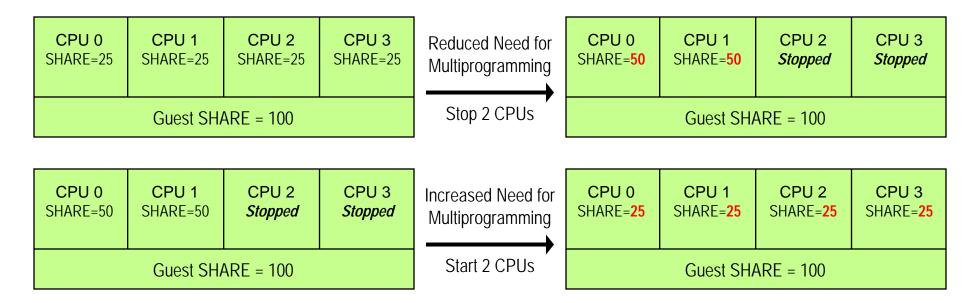
z/VM 5.4: Dynamic Memory Upgrade

- Aka Dynamic Storage Reconfiguration
- Dynamically add real central storage
 - No removal; no expanded storage
- SET STORAGE command response may be delayed
 - Request for small increase in comparison to memory available to partition
- Memory is not all initialized instantaneously
- Also virtualizes for guests that support it
- Some monitor event record anomalies exist
 - Not surfaced in Performance Toolkit
 - Query commands are correct
 - Fixed in service stream with VM64483
- See the z/VM Performance Report



z/VM 5.4: Virtual CPU SHARE Redistribution

- Allows z/VM guests to expand or contract the number of virtual processors it uses without affecting the overall CPU capacity it is allowed to consume
 - Guests can dynamically optimize their multiprogramming capacity based on workload demand
 - Starting and stopping virtual CPUs does not affect the total amount of CPU capacity the guest is authorized to use
 - Linux CPU hotplug daemon starts and stops virtual CPUs based on Linux load average value
- Helps enhance the overall efficiency of a Linux-on-z/VM environment
- Previously, stopped virtual processors were given a portion of the guest share.



Note: Overall CPU capacity for a guest system can be dynamically adjusted using the SHARE setting

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z/VM 5.4: DCSSs Above 2 GB

- DCSSs can be loaded above 2 GB
 - Each DCSS is still limited to 2047 MB
 - Continue to allow multiples, thus aggregate > 2 GB

Requires additional Linux support:

- To use a segment above 2 GB
- To let Linux use 'stacked' contiguous DCSSs as one large block device

Benefits of sharing much larger amounts of memory

- Varies significantly based on configuration and workload.
- Example: throughput improvements of ~35% for static web serving observed with prototype Linux code.

See the z/VM Performance Report



z/VM 5.4: DCSS Loading Time & Tradeoffs

- Support for DCSSs above 2 GB means great amount of memory can be used for DCSSs.
- Performs well once DCSS is built and data saved.
- Placing at higher than necessary addresses results in more memory required for data structures in guest and in z/VM.
- Time required to save or fill a DCSS is non-trivial.
- Two primary methods of defining segments for use by Linux with unique pros and cons
 - SR: Shared read-only access.
 - SN: Shared read/write access, no data saved.



z/VM 5.4: DCSS Tradeoffs, SN vs. SR

DCSS Attribute	SN: Shared R/W; not saved ²	SR: Shared R/O
Non-Volatile		
(Saved to spool)	NO	YES
Initial elapsed time to 'make ready' the file system ¹	Tends to be faster (not necessarily writing to DASD)	Tends to be slower (must write to DASD)
Spool processing for DCSS can delay other spool activity	NO	YES

¹ Making ready the file system involves various steps of defining, copying data into the DCSS, and saving as necessary.

² The DCSS itself is read-write. After data is loaded into the DCSSs, one mounts the file system read-only.



z/VM 5.4: TCP/IP Ethernet Mode

- TCP/IP stack can now run a real OSA in Ethernet mode (aka "layer 2" mode)
- It can also [of course] run a QDIO virtual NIC in Ethernet mode
 - ... thereby letting it couple to an Ethernet-mode guest LAN or Ethernet-mode VSWITCH
- Key findings: for VSWITCH case, according to workload,
 - 0% to 13% improvement in throughput
 - 0% to 7% decrease in CPU time per unit of work
- See the z/VM Performance Report



z/VM 5.4: Telnet IPv6

Our Telnet server can now support IPv6 clients

- New Pascal APIs for IPv6
- Telnet server calls only these new IPv6 APIs
 - Uses IPv6-mapped IPv4 addresses if necessary

Key findings:

- IPv4 regression: impact is "around zero"
- From IPv4 to IPv6:
 - 12% to 23% increase in throughput
 - 3% to 13% decrease in CPU per unit of work

See the z/VM Performance Report



z/VM 5.4: HMC Linux Install

The HMC DVD can now be used to install Linux

- z/VM FTP server will communicate with it
- Alternative to LAN-based server
- Elapsed time for install is much greater than LAN-based server install
 - 11 to 12 times slower
 - Up to 3 hours to do base install
- Do apply the service from the GA RSU
 - APAR PK69228



z/VM 5.4: Upper DAT Structures Above 2 GB

- z/VM 5.3 let page tables (PGMBKs) reside above 2 GB
- z/VM 5.4 further lets segment tables and region tables reside above 2 GB
 - These structures tend to require multiple contiguous frames.
 - Helps reduce fragmentation and long searches
 - Resulting in better performance for some configurations
 - Prereq for increasing supported real memory (which we did not do this release)

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z/VM 5.4: CMS Based SSL Server

Previously, z/VM SSL Server was based on Linux

- Required customers to install Linux in a virtual machine
- Now a CMS based server is used instead of Linux
- APARs shipped in December for this support:
 - PK65850 & PK73085

Performance

- Now exploits the CP Assist for Cryptographic Functions (CPACF) feature where available.
- Very sensitive to the "maximum session" parameter. Do not set this higher than necessary.
- Future performance enhancements are planned.



z/VM 5.4: Other Changes

- PLDV push-through stack: improved efficiency in dequeueing from PLDVs
- Virtual CTC: buffering changes help workloads that do > 32 KB writes to VCTCs
- VSWITCH changes: improved dispatching, suppressed unnecessary load rebalance calculations, increased packet queuing limits
- MDC 8 GB: once MDC reached 8 GB, it stopped doing inserts. Not anymore. ②
- Contiguous available lists: moved low-water and high-water marks closer together – seemed like the right thing for the workloads we tried

VM64439: CMM–VMRM Change

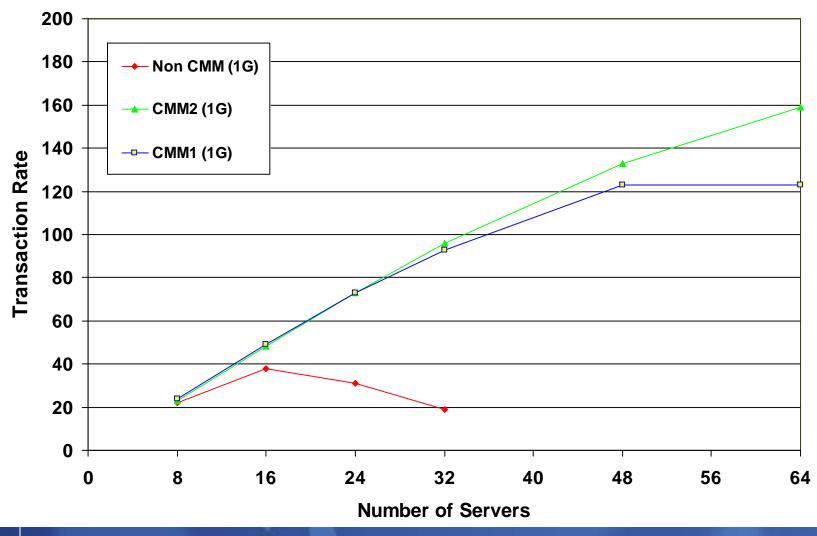
- CMM1 aka CMM & VMRM aka ballooning aka Cooperative Memory Management
- CMM2 aka MEMASSIST aka CMMA aka Collaborative Memory Management
- New CMM-VMRM 64 MB "safety net"
 - In base of z/VM 5.4.0
 - As APAR VM64439 for z/VM 5.2.0 & z/VM 5.3.0
 - z/VM 5.2.0 PTF UM32427
 - z/VM 5.3.0 PTF UM32428

For additional details see: http://www.vm.ibm.com/perf/reports/zvm/html/530cmm.html



Transaction Rate vs. Number of Servers for various Storage Management Products using Apache servers with a virtual storage

size as shown in parentheses in the legend; z9 6 GB / 2 GB

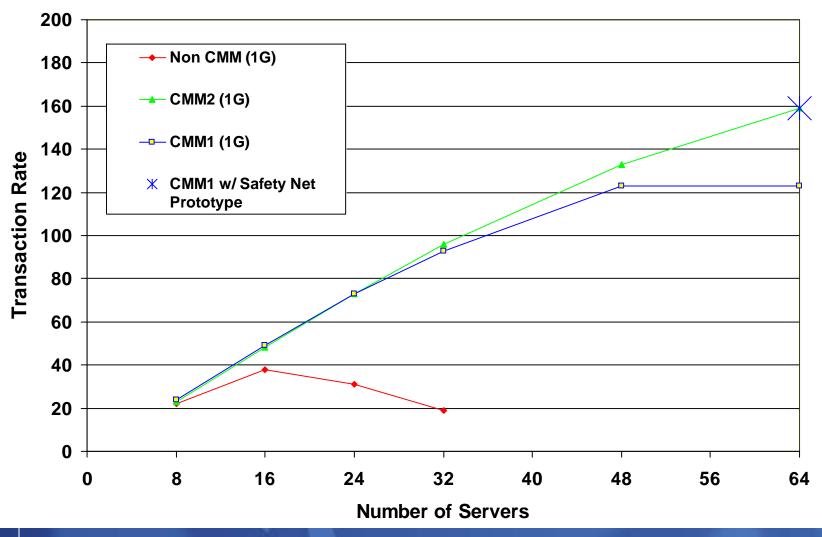


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Transaction Rate vs. Number of Servers

for various Storage Management Products using Apache servers with a virtual storage size as shown in parentheses in the legend



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VM64082 and VM64510: MDC Changes

VM64082

- (z/VM 5.2.0 & 5.3.0). Not on any RSU.
- Rolled into base of z/VM 5.4.0
- This change improves MDC Arbiter choices in some environments (doesn't overbias for MDC).
- Also, makes other environments worse by biasing too much against MDC.
- Use SET MDC command to bound the range of memory to avoid problems.
- VM64510: VM64082 had problems on z/VM 5.3 and 5.4
 - Caused too-frequent steal of pages from MDC
 - VM64510 fixes it



z/VM 5.4: Long Eligible List Stays

- Eligible list is where the scheduler keeps virtual machines to prevent them from running and causing thrashing.
- SET SRM and SET QUICKDSP tuning options can be used to keep the proper guests from being trapped in eligible list.
- There have been known cases of virtual machines trapped in the eligible list for longer than acceptable times in past releases.
- There is evidence that in z/VM 5.4 there is a greater tendency for virtual machines to get trapped when system is not properly tuned.



z/VM: Reorder Processing & Large Resident Counts

- Page reorder is the process of managing user frame owned lists as input to demand scan processing.
 - It includes resetting the HW reference bit.
 - Serializes the virtual machine (all virtual processors).
 - In all releases of z/VM
- It is done periodically on a virtual machine basis.
- The cost of reorder is proportional to the number of <u>resident</u> frames for the virtual machine.
 - Roughly 130 ms/GB resident
 - Delays of ~1 second for guest having 8 GB resident
- Development is investigating improvements



z/VM: VMDUMP Processing

- VMDUMP is a very helpful command for problem determination.
- Some weaknesses:
 - Does not scale well, can take up to 40 minutes per GB.
 - It is not interruptible
- Linux provides a disk dump utility which is much faster relative to VMDUMP.
 - It is disruptive
 - Does not include segments outside the normal virtual machine.



Hardware Concerns

- ECPMF and FCP chpids
- LPAR monitor records
- z10 performance compared to z9



HW: Monitor – Some FCP channels have incorrect data

- Domain 0 Record 20 reports Extended Channel Path Measurement Facility information
- This data appears incorrect for <u>some</u> FCP channels.
- Appears to be a problem with ECPMF, not z/VM
- Investigation continues
- Reference z/VM internal problem number FC01043



HW: LPAR Monitor Records, z10 + z/VM 5.2

- z/VM 5.2 customers running on z10 will not get LPAR monitor records for other partitions on the CEC.
- z/VM 5.3 and z/VM 5.4 are fine.
- Other processors are fine.
- A hardware problem ODT S5619 is open to correct this.
- Corrected with:
 - EC Level: F85901
 - MCL No: 006
 - Driver: D73G



HW: z10 Performance

- Processor cycle time greatly improved over z9
 - ~2.6 times faster (4.4 GHz)
 - Comparable to other platforms
- Laws of physics must be obeyed
- Tradeoffs made in order to achieve above
 - Memory differences
 - Key ops
- ITR ratios (examples see LSPR for most current numbers)
 - z/OS: z10 EC 701 up to **1.62 times** that of the z9 EC 701
 - LSPR z/VM measurements: 1.30 to 1.60
 - z/VM Endicott lab measurements: 1.23 to 2.05

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HW: z10 Performance Depends On....

Number of processors

Fewer processors, better ITRR

Storage references

- Smaller memory footprints, better ITRR
- Data movement
 - Less data movement, better ITRR
- Virtual I/O to real devices
 - Less virtual I/O, better ITRR
- Storage overcommitment
 - Less over commitment, better ITRR
- Amount of memory involved in long searches
 - Shorter & less frequent searches, better ITRR
- Exploitation of new features
 - More exploitation of features, better ITRR



HW: z10 Performance: Setting Proper Expectations

- z10 is a great machine, with a number of excellent attributes.
- Care must be taken when sizing migrations from z9 to z10.
- Additional Information:
 - -LSPR Q & A (complete)
 - Discuss range and factors affecting
 - Pointer to z/VM web page
 - -z/VM web page
 - <u>http://www.vm.ibm.com/perf/z10.html</u>
 - "To MIPS or Not to MIPS, That is the Question!" by Gary King
 - <u>http://shareew.prod.web.sba.com/proceedingmod/abstract.cfm?abstract_id=17583</u>



Summary

- Specialty engines and DMU help z/VM keep up with hardware's capabilities
- Several other enhancements in z/VM 5.4
- z10 is a good machine, but attention to detail is required
- Check out the z/VM performance report: http://www.vm.ibm.com/perf/reports/zvm/html/

- z/VM 5.4 topics published on September 19, 2008